5.2 Algorithm conventions

We often use a numbered list to specify steps in an algorithm. These algorithms are used to clarify semantics. In practice, there may be more efficient algorithms available to implement a given feature.

When an algorithm is to produce a value as a result, we use the directive *return x* to indicate that the result of the algorithm is the value of x and that the algorithm should terminate. We use the notation Result(n) as shorthand for *the result of step n*. We also use Type(x) as shorthand for *the type of x*.

Mathematical operations such as addition, subtraction, negation, multiplication, division, and the mathematical functions defined later in this section should always be understood as computing exact mathematical results on mathematical real numbers, which do not include infinities and do not include a negative zero that is distinguished from positive zero. Algorithms in this standard that model floating-point arithmetic include explicit steps, where necessary, to handle infinities and signed zero and to perform rounding. If a mathematical operation or function is applied to a floating-point number, it should be understood as being applied to the exact mathematical value represented by that floating-point number. Such a floating-point number must be finite, and if it is **+0** or **−0** then the corresponding mathematical value is simply 0.

The mathematical function abs(x) yields the absolute value of x, which is $-x$ if x is negative (less than zero) and otherwise is x itself.

The mathematical function sign(x) yields 1 if x is positive and –1 if x is negative. The sign function is not used in this standard for cases when x is zero.

The notation "x modulo y" (y must be finite and nonzero) computes a value k of the same sign as y such that abs(k) < abs(y) and $x-k = q \cdot y$ for some integer q.

The mathematical function floor(x) yields the largest integer (closest to positive infinity) that is not larger than x. Note that floor(x) = x–(x modulo 1).

If an algorithm is defined to *generate a runtime error*, execution of the algorithm is terminated and no result is returned. The calling algorithms are also terminated, until an algorithm step is reached that explicitly deals with the error. The same applies for exceptions that are explicitly thrown. See section 12.1. The algorithm step that deals with the runtime error, or the explicitly thrown exception, has available to it the details about the error, or the value thrown by the `throw` statement, respectively.

8.9 The Completion Type

The internal Completion type is not a language data type. It is defined by this specification purely for expository purposes. An implementation of ECMAScript must behave as if it produced and operated upon Completion values in the manner described here. However, a value of the Completion type is used only as an intermediate result of statement evaluation and cannot be stored as the value of a variable or property.

The Completion type is used to explain the behavior of statements (`break`, `continue`, `return` and `throw`) that perform nonlocal transfers of control. Values of the Completion type are triples of the form (*type*, *value*, *target*), where *type* is one of **normal**, **break**, **continue**, **return**, or **throw**, *value* is any ECMAScript value, or **empty**, and *target* is any ECMAScript identifier, or **empty**.

The term "abrupt completion" refers to any completion with a reason value other than **normal**.

Invoking the [[Call]] or [[Construct]] method of a Function object, amounts to the evaluation of a *Block* (see section 12.1) in an appropriate Execution Context (see section 10). The result of evaluating a *Block* is of the Completion Type. This value should not be returned as the result of the method invocation, or it might end up being stored in a variable or property. Instead, the *value* field of the completion value becomes the result of the invocation, except that an **empty** value is replaced with **undefined**. If the completion value is of *type* **throw**, execution of the algorithm that invoked the method should proceed as if a runtime error has occurred, see section 5.2.

12 Statements

Syntax

Statement **:**
Block
FunctionDeclaration
VariableStatement
EmptyStatement
ExpressionStatement
IfStatement
IterationStatement
ContinueStatement
BreakStatement
ReturnStatement
WithStatement
LabeledStatement
SwitchStatement
ThrowStatement
TryStatement

Semantics

A *Statement* can be part of a *LabeledStatement*, which itself can be part of a *LabeledStatement*, and so on. The labels introduced this way are collectively referred to as the "current label set" when describing the semantics of individual statements. A *LabeledStatement* has no semantic meaning other than the introduction of a label to a *label set*. An *IterationStatement*, or *SwitchStatement* that is not part of a *LabeledStatement* is regarded as possessing a label set containing a single element, **empty**.

12.1 Block

Syntax

Block **:**
{ *StatementList*$_{opt}$ **}**

StatementList **:**
Statement
StatementList Statement

Semantics

The production *Block* **: { }** is evaluated as follows:

1. Return (normal, **empty**, **empty**).

The production *Block* **: {** *StatementList* **}**is evaluated as follows:

1. Evaluate *StatementList*.
2. Return Result(1).

The production *StatementList* **:** *Statement* is evaluated as follows:

1. Evaluate *Statement*.
2. If an exception value was thrown during the evaluation of *Statement*, go to step 7.
3. If a runtime error occurred during the evaluation of *Statement*, go to step 5.
4. Return Result(1).
5. Construct an appropriate Error object.
6. Return (**throw**, Result(5), **empty**).
7. Return (**throw**, *V*, **empty**) where *V* is the exception value thrown during the evaluation of *Statement*.

The production *StatementList* **:** *StatementList Statement* is evaluated as follows:

1. Evaluate *StatementList*.
2. If Result(1).*type* = **break** and Result(1).*target* occurs in the current label set**,** return (**normal**, Result(1).*value*, **empty**).

3. If Result(1) is an abrupt completion, return Result(1).
4. Evaluate *Statement*.
5. If Result(4).*value* = **empty**, let *V* = Result(1).*value*, otherwise let *V* = Result(4).*value*.
6. If Result(4).*type* = **break** and Result(4).*target* occurs in the current label set**,** return (**normal**, Result(4).*value*, **empty**).
7. Return (Result(4).*type*, *V*, Result(4).*target*).

12.2 Variable statement

Syntax

VariableStatement **:**
  `var` *VariableDeclarationList* `;`

VariableDeclarationList **:**
  VariableDeclaration
  VariableDeclarationList `,` *VariableDeclaration*

VariableDeclaration **:**
  Identifier Initializer$_{opt}$

Initializer **:**
  `=` *AssignmentExpression*

Description

If the variable statement occurs inside a *FunctionDeclaration*, the variables are defined with function-local scope in that function, as described in section **Error! Reference source not found.**. Otherwise, they are defined with global scope, that is, they are created as members of the global object, as described in section **Error! Reference source not found.**. Variables are created when the execution scope is entered. A *Block* does not define a new execution scope. Only *Program* and *FunctionDeclaration* produce a new scope. Variables are initialized to the `undefined` value when created. A variable with an *Initializer* is assigned the value of its *AssignmentExpression* when the *VariableStatement* is executed, not when the variable is created.

Semantics

The production *VariableStatement* **:** `var` *VariableDeclarationList* `;` is evaluated as follows:

1. Evaluate *VariableDeclarationList*.
2. Return (**normal**, **empty**, **empty**).

The production *VariableDeclarationList* **:** *VariableDeclaration* is evaluated as follows:

1. Evaluate *VariableDeclaration*.

The production *VariableDeclarationList* **:** *VariableDeclarationList* `,` *VariableDeclaration* is evaluated as follows:

1. Evaluate *VariableDeclarationList*.
2. Evaluate *VariableDeclaration*.

The production *VariableDeclaration* **:** *Identifier* is evaluated as follows:

1. Return a string value containing the same sequence of characters as in the *Identifier*.

The production *VariableDeclaration* **:** *Identifier Initializer* is evaluated as follows:

1. Evaluate *Identifier*.
2. Evaluate *Initializer*.
3. Call GetValue(Result(2)).
4. Call PutValue(Result(1), Result(3)).
5. Return a string value containing the same sequence of characters as in the *Identifier*.

The production *Initializer* **:** `=` *AssignmentExpression* is evaluated as follows:

1. Evaluate *AssignmentExpression*.
2. Return Result(1).

12.3 Empty statement

Syntax

EmptyStatement **:**
  `;`

Semantics

The production *EmptyStatement* **:** `;` is evaluated as follows:

1. Return (**normal**, **empty**, **empty**).

12.4 Expression statement

Syntax

ExpressionStatement **:**
  Expression `;`

Semantics

The production *ExpressionStatement* **:** *Expression* `;` is evaluated as follows:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Return (**normal**, Result(2), **empty**).

12.5 The `IF` statement

Syntax

IfStatement **:**
  `if (` *Expression* `)` *Statement* `else` *Statement*
  `if (` *Expression* `)` *Statement*

Each `else` for which the choice of associated `if` is ambiguous shall be associated with the nearest possible **if** that would otherwise have no corresponding `else`.

Semantics

The production *IfStatement* **:** `if (` *Expression* `)` *Statement* `else` *Statement* is evaluated as follows:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Call ToBoolean(Result(2)).
4. If Result(3) is **false**, go to step 8.
5. Evaluate the first *Statement*.
6. If Result(5).*type* = **break** and Result(5).*target* occurs in the current label set**,** return (**normal**, Result(5).*value*, **empty**).
7. Return Result(5).
8. Evaluate the second *Statement*.
9. If Result(8).*type* = **break** and Result(8).*target* occurs in the current label set**,** return (**normal**, Result(8).*value*, **empty**).
10. Return Result(8).

The production *IfStatement* **:** `if (` *Expression* `)` *Statement* is evaluated as follows:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Call ToBoolean(Result(2)).
4. If Result(3) is **false**, return (**normal**, **empty**, **empty**).

5. Evaluate *Statement*.
6. If Result(5).*type* = **break** and Result(5).*target* occurs in the current label set**,** return (**normal**, Result(5).*value*, **empty**).
7. Return Result(5).

12.6 Iteration statements

Syntax

IterationStatement **:**
 `do` *Statement* **`while (`** *Expression* **`);`**
 `while (` *Expression* **`)`** *Statement*
 `for (` $Expression_{opt}$ **`;`** $Expression_{opt}$ **`;`** $Expression_{opt}$ **`)`** *Statement*
 `for ( var` *VariableDeclarationList* **`;`** $Expression_{opt}$ **`;`** $Expression_{opt}$ **`)`** *Statement*
 `for (` *LeftHandSideExpression* **`in`** *Expression* **`)`** *Statement*
 `for ( var` *Identifier* $Initializer_{opt}$ **`in`** *Expression* **`)`** *Statement*

12.6.1 The `do…while` Statement

The production **`do`** *Statement* **`while (`** *Expression* **`);`** is evaluated as follows:

1. Let *V* = **empty**.
2. Evaluate *Statement*.
3. If Result(2).*value* is not **empty**, let *V* = Result(2).*value*.
4. If Result(2).*type* = **continue** and Result(2).*target* is in the current label set, go to 2.
5. If Result(2).*type* = **break** and Result(2).*target* is in the current label set, return (**normal**, *V*, **empty**).
6. If Result(2) is an abrupt completion, return Result(2).
7. Evaluate Expression.
8. Call GetValue(Result(7)).
9. Call ToBoolean(Result(8)).
10. If Result(9) is true, go to step 2.
11. Return (**normal**, *V*, **empty**);

12.6.2 The `while` statement

The production *IterationStatement* **:** **`while (`** *Expression* **`)`** *Statement* is evaluated as follows:

1. Let *V* = **empty**.
2. Evaluate *Expression*.
3. Call GetValue(Result(2)).
4. Call ToBoolean(Result(3)).
5. If Result(4) is **false**, return (**normal**, *V*, **empty**).
6. Evaluate *Statement*.
7. If Result(6).*value* is not **empty**, let *V* = Result(6).*value*.
8. If Result(6).*type* = **continue** and Result(6).*target* is in the current label set, go to 2.
9. If Result(6).*type* = **break** and Result(6).*target* is in the current label set, return (**normal**, *V*, **empty**).
10. If Result(6) is an abrupt completion, return Result(6).
11. Go to step 2.

12.6.3 The `for` statement

The production *IterationStatement* **:** **`for (`** $Expression_{opt}$ **`;`** $Expression_{opt}$ **`;`** $Expression_{opt}$ **`)`** *Statement* is evaluated as follows:

1. If the first *Expression* is not present, go to step 4.
2. Evaluate the first *Expression*.
3. Call GetValue(Result(2)). (This value is not used.)
4. Let *V* = **empty**.
5. If the second *Expression* is not present, go to step 10.
6. Evaluate the second *Expression*.

7. Call GetValue(Result(6)).
8. Call ToBoolean(Result(7)).
9. If Result(8) is **false**, go to step 19.
10. Evaluate *Statement*.
11. If Result(10).*value* is not **empty**, let *V* = Result(10).*value*
12. If Result(10).*type* = **break** and Result(10).*target* is in the current label set, go to step 19.
13. If Result(10).*type* = **continue** and Result(10).*target* is in the current label set, go to step 15..
14. If Result(10) is an abrupt completion, return Result(10).
15. If the third *Expression* is not present, go to step 5.
16. Evaluate the third *Expressio*n.
17. Call GetValue(Result(16). (This value is not used.)
18. Go to step 5.
19. Return (**normal**, *V*, **empty**).

The production *IterationStatement* **: `for` `(` `var`** *VariableDeclarationList* **`;`** *Expression*$_{opt}$ **`;`** *Expression*$_{opt}$ **`)`** *Statement* is evaluated as follows:

1. Evaluate *VariableDeclarationList*.
2. Let *V* = **empty**.
3. If the second *Expression* is not present, go to step 8.
4. Evaluate the second *Expression.*
5. Call GetValue(Result(4)).
6. Call ToBoolean(Result(5)).
7. If Result(6) is **false**, go to step 15.
8. Evaluate *Statement*.
9. If Result(8).*value* is not **empty**, let *V* = Result(8).*value*.
10. If Result(8).*type* = **break** and Result(8).*target* is in the current label set, go to step 17.
11. If Result(8).*type* = **continue** and Result(8).*target* is in the current label set, go to step 13.
12. If Result(8) is an abrupt completion, return Result(8).
13. If the third *Expression* is not present, go to step 3.
14. Evaluate the third *Expression*.
15. Call GetValue(Result(14)). (This value is not used.)
16. Go to step 3.
17. Return (**normal**, *V*, **empty**).

12.6.4 The `for..in` statement

The production *IterationStatement* **: `for` `(`** *LeftHandSideExpression* **`in`** *Expression* **`)`** *Statement* is evaluated as follows:

1. Evaluate the *Expression*.
2. Call GetValue(Result(1)).
3. Call ToObject(Result(2)).
4. Let *V* = **empty**.
5. Get the name of the next property of Result(3) that doesn't have the DontEnum attribute. If there is no such property, go to step 14.
6. Evaluate the *LeftHandSideExpression* (it may be evaluated repeatedly).
7. Call PutValue(Result(6), Result(5)).
8. Evaluate *Statement*.
9. If Result(8).*value* is not **empty**, let *V* = Result(8).*value*.
10. If Result(8).*type* = **break** and Result(8).*target* is in the current label set, go to step 14.
11. If Result(8).*type* = **continue** and Result(8).*target* is in the current label set, go to step 5.
12. If Result(8) is an abrupt completion, return Result(8).
13. Go to step 5.
14. Return (**normal**, *V*, **empty**).

The production *IterationStatement* **: `for` `(` `var`** *VariableDeclaration* **`in`** *Expression* **`)`** *Statement* is evaluated as follows:

1. Evaluate *VariableDeclaration*.
2. Evaluate *Expression*.
3. Call GetValue(Result(2)).
4. Call ToObject(Result(3)).
5. Let *V* = **empty**.
6. Get the name of the next property of Result(4) that doesn't have the DontEnum attribute. If there is no such property, go to step 19.
7. Evaluate Result(1) as if it were an *Identifier*; see **Error! Reference source not found.** (yes, it may be evaluated repeatedly).
8. Call PutValue(Result(7), Result(6)).
9. Evaluate *Statement*.
10. If Result(9).*value* is not **empty**, let *V* = Result(9).*value*.
11. If Result(9).*type* = **break** and Result(9).*target* is in the current label set, go to step 15.
12. If Result(9).*type* = **continue** and Result(9).*target* is in the current label set, go to step 6.
13. If Result(8) is an abrupt completion, return Result(8).
14. Go to step 6.
15. Return (**normal**, *V*, **empty**).

The mechanics of enumerating the properties (step 5 in the first algorithm, step 6 in the second) is implementation dependent. The order of enumeration is defined by the object. Properties of the object being enumerated may be deleted during enumeration. If a property that has not yet been visited during enumeration is deleted, then it will not be visited. If new properties are added to the object being enumerated during enumeration, the newly added properties are not guaranteed to be visited in the active enumeration.

Enumerating the properties of an object includes enumerating properties of its prototype, and the prototype of the prototype, and so on, recursively; but a property of a prototype is not enumerated if it is "shadowed" because some previous object in the prototype chain has a property with the same name.

12.7 The `CONTINUE` statement

Syntax

ContinueStatement **:**
`continue` [no *LineTerminator* here] $Identifier_{opt}$ `;`

Semantics

A program is considered syntactically incorrect if either of the following are true:

- The program contains a `continue` statement without the optional *Identifier*, which is not nested, directly or indirectly, within an *IterationStatement*.
- The program contains a `continue` statement with the optional *Identifier*, where *Identifier* does not appear in the label set of an enclosing *IterationStatement*.

A *ContinueStatement* without an *Identifier* is evaluated as follows:

1. Return (**continue**, **empty**, **empty**).

A `continue` statement with the optional *Identifier* is evaluated as follows:

1. Return (**continue**, **empty**, *Identifier*).

12.8 The `BREAK` statement

Syntax

BreakStatement **:**
`break` [no *LineTerminator* here] $Identifier_{opt}$ `;`

Semantics

A program is considered syntactically incorrect if either of the following are true:

- The program contains a **`break`** statement without the optional *Identifier*, which is not nested, directly or indirectly, within an *IterationStatement* or a *SwitchStatement*.
- The program contains a **`break`** statement with the optional *Identifier*, where *Identifier* does not appear in the label set of an enclosing *Statement*.

A *BreakStatement* without an *Identifier* is evaluated as follows:

1. Return (**break**, **empty**, **empty**).

A **`break`** statement with an *Identifier* is evaluated as follows:

1. Return (**break**, **empty**, *Identifier*).

12.9 The `RETURN` statement

Syntax

ReturnStatement **:**
  `return` [no *LineTerminator* here] *Expression$_{opt}$* **`;`**

Semantics

An ECMAScript program is considered syntactically incorrect if it contains a **`return`** statement that is not within the *Block* of a *FunctionDeclaration*. It causes a function to cease execution and return a value to the caller. If *Expression* is omitted, the return value is the **`undefined`** value. Otherwise, the return value is the value of *Expression*.

The production *ReturnStatement* **::** **`return`** [no *LineTerminator* here] *Expression$_{opt}$* **`;`** is evaluated as:

1. If the *Expression* is not present, return (**return**, **`undefined`**, **empty**).
2. Evaluate *Expression*.
3. Call GetValue(Result(2)).
4. Return (**return**, Result(3), **empty**).

12.10 The `WITH` statement

Syntax

WithStatement **:**
  `with (` *Expression* **`)`** *Statement*

Description

The **`with`** statement adds a computed object to the front of the scope chain of the current execution context, then executes a statement with this augmented scope chain, then restores the scope chain.

Semantics

The production *WithStatement* **:** **`with (`** *Expression* **`)`** *Statement* is evaluated as follows:

1. Evaluate Expression.
2. Call GetValue(Result(1)).
3. Call ToObject(Result(2)).
4. Add Result(3) to the front of the scope chain.
5. Evaluate *Statement* using the augmented scope chain from step 4.
6. Remove Result(3) from the front of the scope chain.
7. Return Result(5).

Discussion

Note that no matter how control leaves the embedded *Statement*, whether normally or by some form of abrupt completion, the start of the scope chain is always restored to its former state.

12.11 The SWITCH Statement

Syntax

SwitchStatement **:**
    `switch (` *Expression* **`)`** *CaseBlock*

CaseBlock **:**
    `{` $CaseClauses_{opt}$ **`}`**
    `{` $CaseClauses_{opt}$ *DefaultClause* $CaseClauses_{opt}$ **`}`**

CaseClauses **:**
    CaseClause
    CaseClauses CaseClause

CaseClause **:**
    `case` *Expression* **`:`** $StatementList_{opt}$

DefaultClause **:**
    `default :` $StatementList_{opt}$

Semantics

The production *SwitchStatement* **:** **`switch (`** *Expression* **`)`** *CaseBlock* is evaluated as follows:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Evaluate *CaseBlock*, passing it Result(2) as a parameter.
4. If Result(3).*type* = **break** and Result(3).*target* is in the current label set, return (**normal**, Result(3).*value*, **empty**).
5. Return Result(3).

The production *CaseBlock* **: `{`** *CaseClauses DefaultClause CaseClauses* **`}`** is given an input parameter, *input*, and is evaluated as follows:

1. Let *A* be the list of *CaseClause* items in the first *CaseClauses*, in source text order.
2. For the next *CaseClause* in *A*, evaluate *CaseClause*. If there is no such *CaseClause*, go to step 7.
3. If *input* is not equal to Result(2), as defined by the !== operator, go to step 2.
4. Evaluate the *StatementList* of this *CaseClause*.
5. If Result(4) is an abrupt completion then return Result(4).
6. Go to step 13.
7. Let *B* be the list of *CaseClause* items in the second *CaseClauses*, in source text order.
8. For the next *CaseClause* in *B*, evaluate *CaseClause*. If there is no such *CaseClause*, go to step 15.
9. If *input* is not equal to Result(8), as defined by the !== operator, go to step 8.
10. Evaluate the *StatementList* of this *CaseClause*.
11. If Result(10) is an abrupt completion then return Result(10).
12. Go to step 18.
13. For the next *CaseClause* in *A*, evaluate the *StatementList* of this *CaseClause*. If there is no such *CaseClause*, go to step 15.
14. If Result(13) is an abrupt completion then return Result(13).
15. Execute the *StatementList* of *DefaultClause*.
16. If Result(15) is an abrupt completion then return Result(15).
17. Let *B* be the list of *CaseClause* items in the second *CaseClauses*, in source text order.
18. For the next *CaseClause* in *B*, evaluate the *StatementList* of this *CaseClause*. If there is no such *CaseClause*, return (**normal**, **empty**, **empty**).
19. If Result(18) is an abrupt completion then return Result(18).
20. Go to step 18.

The production *CaseClause* **: `case`** *Expression* **`:`** $StatementList_{opt}$ is evaluated as follows:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Return Result(2).

Note that evaluating *CaseClause* does not execute the associated *StatementList*. It simply evaluates the *Expression* and returns the value, which the *CaseBlock* algorithm uses to determine which *StatementList* to start executing.

12.12 Labeled Statements

Syntax

LabeledStatement **:**
    Identifier **`:`** *Statement*

Semantics

A *Statement* may be prefixed by a label. Labeled statements are only used in conjunction with labeled **`break`** and **`continue`** statements. ECMAScript has no **`goto`** statement.

An ECMAScript program is considered syntactically incorrect if it contains a *LabeledStatement* that is enclosed by a *LabeledStatement* with the same *Identifier* as label.

The production *Identifier* ***`:`*** *Statement* is evaluated by adding *Identifier* to the label set of *Statement* and then evaluating *Statement*. If the *LabeledStatement* itself has a non-empty label set, these labels are also added to the label set of *Statement* before evaluating it.

Prior to the evaluation of a *LabeledStatement*, the contained *Statement* is regarded as possessing an empty label set, except if it is an *IterationStatement* or a *SwitchStatement*, in which case it is regarded as possessing a label set consisting of the single element, **empty**.

12.10 The `THROW` statement

Syntax

ThrowStatement :
    `throw` [no *LineTerminator* here] *Expression ;*

Semantics

The production *ThrowStatement* **::** ***`throw`*** [no *LineTerminator* here] *Expression* ***`;`*** is evaluated as:

1. Evaluate *Expression*.
2. Call GetValue(Result(1)).
3. Return (**throw**, Result(2), **empty**), behaving as if a runtime error has occurred. See section 5.2.

12.11 The `TRY` statement

Syntax

TryStatement :
    `try` *Block* **`catch (var`** *Identifier* **`)`** *Block*

Description

The **`try`** statement encloses a block of code in which an exceptional condition can occur, such as a runtime error or a **`throw`** statement. The **`catch`** clause provides the exception-handling code. The identifier introduces a local variable that is created when the execution scope containing the try statement is entered.

Semantics

The production *TryStatement* ***::*** ***`try`*** *Block* **`catch (var`** *Identifier* **`)`** *Block* ; is evaluated as follows:

1. Evaluate the first *Block*.
2. If Result(1).*type* is not **throw**, return Result(1).
3. Evaluate *Identifier*.
4. Call PutValue(Result(3), V).

5. Evaluate the second *Block*.
6. Return Result(5).

14 Program

Syntax

Program **:**
    SourceElements

SourceElements **:**
    SourceElement
    SourceElements SourceElement

SourceElement **:**
    Statement
    FunctionDeclaration

The production *Program* **:** *SourceElements* is evaluated as follows:

1. Process *SourceElements* for function declarations.
2. Evaluate *SourceElements*.
3. Return Result(2).

The production *SourceElements***:** *SourceElement* is processed for function declarations as follows:

1. Process *SourceElement* for function declarations.

The production *SourceElements***:** *SourceElement* is evaluated as follows:

1. Evaluate *SourceElement*.
2. Return Result(1).

The production *SourceElements***:** *SourceElements SourceElement* is processed for function declarations as follows:

1. Process *SourceElements* for function declarations.
2. Process *SourceElement* for function declarations.

The production *SourceElements***:** *SourceElements SourceElement* is evaluated as follows:

1. Evaluate *SourceElements*.
2. Evaluate *SourceElement*.
3. If Result(2).*value* = **empty**, let Result(2).*value* = Result(1).*value*.
4. Return Result(2).

The production *SourceElement***:** *Statement* is processed for function declarations by taking no action.

The production *SourceElement***:** *Statement is* evaluated as follows:

1. Evaluate *Statement*.
2. Return Result(1).

The production *SourceElement***:** *FunctionDeclaration* is processed for function declarations as follows:

1. Process *FunctionDeclaration* for function declarations.

The production *SourceElement***:** *FunctionDeclaration is* evaluated as follows:

1. Return (**normal**, **empty**, **empty**).

15.1.2 Function properties of the global object

15.1.2.1 eval(x)

When the **`eval`** function is called with one argument *x*, the following steps are taken:

1. If *x* is not a string value, return *x*.

2. Parse *x* as an ECMAScript *Program*. If the parse fails, generate a runtime error.
3. Evaluate the program from step 2.
4. If Result(3).*type* = **throw**, return Result(3), behaving as if a runtime error has occurred, see section 5.2.
5. If Result(3).*value* is not **empty**, return Result(3).*value*.
6. Return **undefined**.