

[[strawman:  
inherited\_explicit\_soft\_fields]]Trace: » es3.1 » start »  
iterators\_and\_generators

» const\_functions » inherited\_explicit\_soft\_fields

## Explicit Inherited Soft Fields

The following derived abstraction combines the explicitness of **explicit soft own fields** with the visibility across inheritance chains of **inherited soft fields**. Below is an executable specification as a wrapper around **weak\_maps**. This strawman page suggests standardizing this derived abstraction because a primitive implementation is likely to be more efficient than the code below.

As with our previous "EphemeronTable", the name "ExplicitSoftField" is only a placeholder until someone suggests an acceptable name.

### Table of Contents

- Explicit Inherited Soft Fields
  - A Less Aggressive GC Contract?
    - Why it might not matter
    - What if it does matter?
- See

```
const ExplicitSoftField() {
  const et = WeakMap();
  const mascot = {}; // fresh and encapsulated, thus differs from any possible
  provided value.
  return Object.freeze({
    get: const(base) {
      while (base !== null) {
        const result = et.get(base);
        if (result !== undefined) {
          return result === mascot ? undefined : result;
        }
        base = Object.getPrototypeOf(base);
      }
      return undefined;
    },
    set: const(key, val) {
      et.set(key, val === undefined ? mascot : val);
    },
    has: const(key) {
      return et.get(key) !== undefined;
    },
    delete: const(key) {
      et.set(key, undefined);
    }
  });
}
```

## A Less Aggressive GC Contract?

At (es-discuss:011705) I (MarkM) wrote regarding the contrast with the **names** strawman:

If the only semantic difference is (not normally observable) less aggressive GC obligations, great. I'm confident we can converge those.

Given the **gc semantics** of **weak maps**, the gc obligations implied by the above `ExplicitSoftField` implementation vs. **names** seems to be:

Labels	ExplicitSoftFields	Names
-	Given field $F$ where $F.get(K) === V$	Given name $F$ where $K[F] === V$
-	Similarities	
a	F and K may retain V	
b	F must not retain K	
c	F without K must not retain V	
-	Differences	
x	K must not retain F	K may retain F
y	K without F must not retain V	K without F may retain V

## Why it might not matter

The **classes as sugar** strawman uses `ExplicitSoftFields` to implement class-private instance variables. Similar classes strawmen have suggested using **names** for similar purposes. How do the above differences affect the gc semantics of these class proposals?

For a given class  $C$ , let's call  $IC$  the set of all instances of  $C$ . The `ExplicitSoftField` or name  $F$  created by the desugaring is reachable only from  $C$  and from  $IC$ , and cannot escape from this set. The keys of  $F$  are exactly all the members of  $IC$ . Under these circumstances, the implied GC obligations seem to be exactly the same. To see why, say that some non-empty subset of  $IC$ ,  $SIC$ , is reachable from outside  $C$  and  $IC$ . Since  $F$  is reachable from all members of  $SIC$  (independent of whether  $F$  is an `ExplicitSoftField` or name),  $SIC$ 's non-emptiness implies that  $F$  may be retained. By rule #b,  $SIC$  and  $F$  and  $C$  together must not retain the remaining members of  $IC$  outside of  $SIC$ . Since these are not retained, by rule #c the private facets of these instances are also not retained. For all members of  $SIC$ , since they retain  $F$ , by rule #a,  $SIC$  together with  $F$  retain all the private facets of  $SIC$ .

The reason that the different obligations of `ExplicitSoftFields` and names don't matter for classes is that there are no circumstances where a  $K$  may be reachable but the corresponding  $F$  might not be, since the  $K$ s in question – the instances of  $C$  – reference  $F$  anyway.

## What if it does matter?

Nevertheless, we may want to reduce the gc obligations of `ExplicitSoftFields` towards that of names. I'm not sure, but I think only difference #y matters. Rule #x *by itself* would only affect how many empty `ExplicitSoftFields` are retained, since the number of Key-to-Value associations retained is determined by the other rules. We can change our executable spec above to represent these reduced obligations as follows:

```
const ExplicitSoftField = (const() {
  const globalET = WeakMap(); // necessarily reachable
  return const() {
    const et = WeakMap();
    const mascot = {};
    return Object.freeze({
      //...other methods same as before...
      set: const(key, val) {
        et.set(key, val === undefined ? mascot : val); // as before
        globalET.set(key, et);
      }
    });
  };
});
```

## See

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The thread beginning at [WeakMap API questions?](#)

strawman/inherited\_explicit\_soft\_fields.txt · Last modified: 2010/09/06 04:20 by markm

