Ecma/TC39-TG1/2006/022



Minutes of the: held in: on:

Ecma TC39-TG1 Phone conference 3rd May 2006

Attendees

Meeting time 10am PDT.

On phone:

- Francis Cheng, Adobe Systems
- Jeff Dyer, Adobe Systems
- Gary Grossman, Adobe Systems
- Ed Smith, Adobe Systems
- Pratap Lakshman, Microsoft
- Brendan Eich, Mozilla Foundation
- Graydon Hoare, Mozilla Foundation
- Blake Kaplan, Mozilla Foundation
- Dave Herman, Northeastern University
- Lars Thomas Hansen, Opera Software
- Cormac Flanagan, UC Santa Cruz

Agenda

Note new meeting day of week for phone conferences.

- package semantics
- multiple compilation units
- formal type system
- other hot topics

Discussion

- package semantics
 - o packages declare two namespaces, p.q#public and p.q#internal
 - o import p.q is in part like use namespace p.q#public
 - o expression starting p.q.x is rewritten to p.q#public::x
 - even within package p.q, p.q.x is rewritten, so x must be public
 - internal::x or just x would work unless ambiguity requires full path
- <u>multiple compilation units</u>
 - Dave: what if you have package p.q with an x use but no x def; now add x to p.q after the compiler dealt with the first x use
 - Jeff: package compilation ends at verification or loading
 - Dave: so packages do involve separate compilation units

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- Lars: see multiple compilation units for browser constraints
 - Ed: AS3 ignores redefinition
 - two programs define utilities, pgm1 has A&B, pgm2 has B&C
 - they load in the same runtime, in that order
 - AS3 assumes first B is same as second B
 - Brendan: browsers and ES1-3 of course have writeable function bindings
 - so last one wins
 - can we do better than last-wins for global functions or first-wins for classes and packages?
- Ed: hard to share common utility packages without getting too fine-grained
- Dave: what are use-cases for splitting up packages into multiple pieces?
 - Jeff: Java examples to avoid overlarge files
 - Dave: easy to unify that case at load time
 - Ed: package with hash table and tree
 - pgm1 uses hash table
 - pgm2 uses tree
 - Dave: why are those in the same package? Ok, pick a better example
 - Ed: explored a signature checksum scheme to verify Bs don't conflict
 - Ed: another example: graphing components for charting
 - also accessibility addons to the charting package
 - want accessibility stuff in a separate compilation unit
 - Brendan: first one wins is going to be hard to beat
 - Ed: Java does that within a classloader
 - Lars: anyone-wins is going to break on the web
 - Brendan: yeah, many <script src
 <p>cases are like #include, some are more like block-scoped import
 - Ed: Flash has application domains outside the AS3 language
 - you can create a sub-domain to isolate effects
 - lookups start with super-domain then go to sub-domain
 - Pratap: CLR2 has app domains too
 - Gary: Flash took inspiration from that, similar
 - Dave: shadowing is not mutability
 - Brendan: browsers name modules by URI, so no subversion via shadowing
 - Ed: packages are namespaces are URIs, so do tie into security and http caching
 - Graydon: content hashing better than relying on DNS
- <u>formal type system</u> questions
 - String to Boolean
 - no controversy on if, while, for, &&, ||, ! converting
 - var x : Boolean = "hi"
 leff: that converts in A
 - Jeff: that converts in AS3 in bang or tilde
 - compatibility requires this without type annotations
 - Brendan: could be stricter
 - Jeff: refactoring hazard
 - Cormac: tradeoff between type-checking and convenience/migration
 - Return from constructors
 - Brendan: ES1-3 allow function constructor to return a different object
 - Ed: class constructor functions cannot return expr; at all
 - but class ctors can return; to bail early
 - Brendan: different from rule in functions
 - Dave: how does type system talk about type of constructor?
 - so could allow constructors to have Void return type
 - and they could even contain return void 0 or whatever
 - Jeff: in AS3, function f():void{...} means ... cannot return expr;
 - Dave: type Void means can't return a value



- Brendan: then need type Undefined too
- Dave: need to review proper tail calls in light of this
- Ed: try this:
 - Void is type, has value undefined
 - f():void implies extra syntax restriction against return expr;
 - but otherwise doesn't affect type-checking, proper tail calls, etc.
 - Dave: concerned about need to name Undefined or Null in unions, etc.
 - Brendan, Ed: need special restriction on return expr; for
 - constructors
 - setters
 - generators
- o with discussion
 - Using annotations and structural types, one can finally state the precise type of the object named in the with head
 - let declarations in body of with work as elsewhere
 - Apart from these orthogonal goods, can we reform with, or banish it?
 - use strict could banish it to a { use dynamic; ... } block
 - does this really help? migration vs. new code, why do people use with?
- o Global object unknowns
 - Brendan: <u>which prototype</u> proposes immutable String, etc.
 - intrinsic::String vs. String wouldn't differ if we adopt that proposal
 - s.intrinsic::charAt(i) would be different from s.charAt(i) for backward compatibility, to support AOP-ish hacking

```
class String {
```

```
...
intrinsic function charAt(i:uint):String {...}
prototype function charAt(i:*):String {...}
}
```

- intrinsic proposal
 - Use intrinsic instead of something like AS3 or ES4 for early binding.